





The Polyhedral Model Beyond Loops Recursion Optimization and Parallelization Through Polyhedral Modeling

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Outline

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2 Proposed Solution: From Recursive Functions to Optimized Loops

- 3 Case Studies
- 4 Conclusion and Perspectives

1 Introduction

Proposed Solution: From Recursive Functions to Optimized Loops

3 Case Studies

4 Conclusion and Perspectives

Motivation

There may be a huge gap between:

- the statements in a program source code
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- the statements in a program source code
- the instructions actually performed by a given processor architecture

Efficient optimizations may be applied as soon as the actual runtime behavior has been discovered

 dedicated to specific control structures & memory access patterns

Inspiration

Apollo

- Captures a polyhedral behavior of loops at runtime
- Applies the polyhedral model

Memory Accesses Behavior at Runtime from statically non-polyhedral loops!



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Objectives

We are interested in recursive functions:

- 1 whose runtime behavior can be modeled as polyhedral loops
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Objectives

- optimizing recursive functions through transformation into affine loops
- extending the scope of polyhedral optimizations to cover recursive functions



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Implementation



The implementation consists of 3 main steps:

 Recursive Control and Memory Behavior Analysis

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- Recursive Control and Memory Behavior Analysis
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- 3 Polyhedral Optimizations



Input: recursive code

Apply classical LLVM optimization passes

- promote memory to register
- simplify CFG
- dead code elimination

Output: optimized LLVM IR & call graph



Input: optimized IR & call graph



Output: direct & indirect recursions



Input: direct & indirect recursions



Output: instrumented recursive code



Input: Trace of the program execution : Basic Block IDs & Memory Addresses

Nested Loop Reconginition (NLR) algorithm applications:

- program behavior modeling for any measured quantity such as memory accesses
- 2 execution trace compressing
- 3 value prediction

(ketterlin & Clauss, GGO 2008)

Output: Affine Loop Model

Recursion to Affine Loop Nest Transformation



Input: Affine loop model

- Extract NLR resulting loop nests structures
- 2 Construct loops in the LLVM IR using:
 - Instrumented basic blocks
 - Interpolated memory addresses

Output: Iterative code with affine loops

Polyhedral Optimizations



Input: Iterative code with affine loops

- Apply LLVM optimization passes
- Use polly LLVM polyhedral optimizer (Grosser et al., PPL 2012)

Output: Optimized recursive code

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Recursive Matrix Multiplication

```
void MatrixMultiplication(int A[N][N], int B[N][N]){
 static int row=0, column=0, index=0;
   if (row >= N)
          return:
        if(column < N){</pre>
          if(index < N){</pre>
           C[row][column]+= A[row][index]*B[index][column];
                 index++:
                 MatrixMultiplication(A, B);
          index=0;
          column++:
          MatrixMultiplication(A, B);
        3
        column=0;
        row++:
        MatrixMultiplication(A, B);
```

Case Studies

Recursive Matrix Multiplication Analysis Results

```
for i\theta = 0 to N-1
 for i1 = 0 to N-1
   for i2 = 0 to N-1
    val MatrixMultiplication::if.then4 //IR basic block
    load // memory read
    val MEM1 + 4*N*i0 + 4*i2 //memory address in terms of loops indices
    ... //repetitive memory access patterns
    load
    val MEM2 + 4*i1 + 4*N*i2 //4 is the size of an integer
    val load
    val MEM3 + 4*N*i0 + 4*i1
    val store // memory write
    val MEM3 + 4*N*i0 + 4*i1
  val MatrixMultiplication::if.end15
 val MatrixMultiplication:: if.end17
for i\theta = \theta to N*N-1
 for i1 = 0 to N-1
  val MatrixMultiplication::if.end17
  val MatrixMultiplication::if.end15
 val MatrixMultiplication:: if.end15
```

Case Studies

Recursive Matrix Multiplication Experimental Results



Serial execution (gcc -O3)

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Heat

The function compstripe involves interesting linear loops

```
void compstripe(register double **new. register double **old. int lb. int ub)
{
 register int a. b. llb. lub:
 llb = (lb == 0) ? 1 : lb:
 lub = (ub == nx) ? nx - 1 : ub;
 for (a=llb; a < lub; a++) {</pre>
  for (b=1; b < ny-1; b++) {</pre>
    new[a][b] = dtdxsg * (old[a+1][b] - 2 * old[a][b] + old[a-1][b])
           + dtdysg * (old[a][b+1] - 2 * old[a][b] + old[a][b-1])
            + old[a][b]:
   }
 for (a=llb; a < lub; a++)</pre>
   new[a][nv-1] = randb(xu + a * dx. t):
 for (a=llb: a < lub: a++)</pre>
   new[a][0] = randa(xu + a * dx, t);
 if (lb == 0) {
   for (b=0: b < nv: b++)
    new[0][b] = randc(yu + b * dy, t);
 3
 if (ub == nx) {
  for (b=0: b < nv: b++)
    new[nx-1][b] = randd(yu + b * dy, t);
 }
```

Heat Analysis Results

```
for iθ = 0 to Number_of_Steps-1
 for i1 = 0 to 14
  for i2 = 0 to 509
    val compstripe::for.body10, MEM1 + 8224*i1 + 8*i2, MEM2 + 8224*i1 + 8*i2, MEM3 + 8224*i1 + 8*i2
      . MEM4 + 8224*i1 + 8*i2 . MEM5 + 8224*i1 + 8*i2. MEM6 + 8224*i1 + 8*i2
 for i1 = 0 to 14
  val compstripe::for.body63. MEM7 + 8224*il
 for i1 = 0 to 14
  val compstripe::for.body81, MEM8 + 8224*i1
 for il = 0 to 511
  val compstripe::for.body97, MEM9 + 8*i1
 for i1 = 0 to 61
  for i2 = 0 to 15
    for i3 = 0 to 509
     val comostripe::for.body10. MEM10 + 131584*i1 + 8224*i2 + 8*i3. MEM11 + 131584*i1 + 8224*i2 + 8*i3. MEM12 + 131584*i1 + 8224*i2 + 8*i3
        . MEM13 + 131584*i1 + 8224*i2 + 8*i3. MEM14 + 131584*i1 + 8224*i2 + 8*i3. MEM15 + 131584*i1 + 8224*i2 + 8*i3
  for i2 = 0 to 15
    val compstripe::for.bodv63 . MEM16 + 131584*i1 + 8224*i2
  for i7 = 0 to 15
    val compstripe::for.body81 , MEM17 + 131584*i1 + 8224*i2
 for i1 = 0 to 14
  for i2 = 0 to 509
    val compstripe::for.body10, MEM18 + 8224*i1 + 8*i2, MEM19 + 8224*i1 + 8*i2, MEM20 + 8224*i1 + 8*i2
      . MEM21 + 8224*i1 + 8*i2. MEM22 + 8224*i1 + 8*i2. MEM23 + 8224*i1 + 8*i2
 for i1 = 0 to 14
  val compstripe::for.body63 . MEM24 + 8224*i1
 for i1 = 0 to 14
  val compstripe::for.body81 , MEM25 + 8224*i1
 for il = 0 to 511
  val compstripe::for.body115 , MEM26 + 8*i1
```

Heat Experimental Results



The codes have been parallelized by Pluto using OpenMP 24 threads (AMD Opteron 6172 2x12-cores - gcc -O3 -fopenmp)

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Conclusion

A proof of concept for an automatic recursion-to-affine-loop transformation:

- involving static and dynamic analysis
- transformation passes
- polyhedral optimizers

Achievements

- extends the polyhedral model applicability to non-loop control structures
- 2 brings the handled recursive functions to a higher level of optimizations

Future Works

Our future works include:

- Performing dynamic analysis for recursive behavior at runtime
- 2 Inducing verification features to obtain a predictive model
- 3 Tackling input dependent recursive codes

Thank you

Questions ?