

# Polyhedral Optimization For JavaScript: The Challenges

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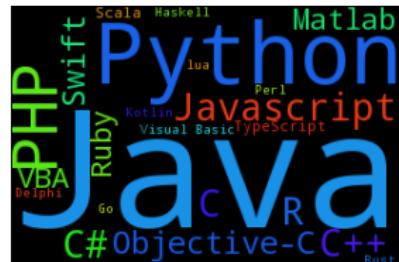
# The JavaScript Language

Created in 1995 at Netscape

- To implement dynamism in web pages
- High level and dynamic
- ECMAScript standard

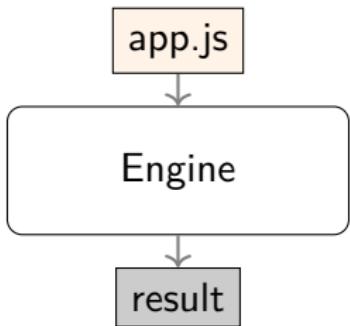
Now widely used both on client and server sides

- For complex applications
- Because of portability
- Because of performances



5th language - PYPL index

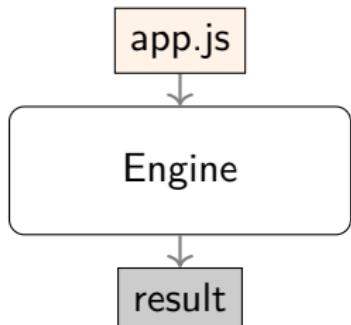
# JavaScript Implementation



## Widespread engines

- SpiderMonkey - *Mozilla*
- JavaScriptCore - *Apple*
- V8 - *Google*
- Chakra - *Microsoft*

# JavaScript Implementation



## Widespread engines

- SpiderMonkey - *Mozilla*
- JavaScriptCore - *Apple*
- V8 - *Google*
- Chakra - *Microsoft*

## Browser wars

- Complex optimization
- Nevertheless, **no parallelism**

# You Said Polyhedral Model And JavaScript Together

## How to cope with dynamism?

- Static Control Parts (SCoPs) cannot be detected statically
- When and how detect polyhedral opportunities?

## Is it worthwhile to use the polyhedral model at runtime?

- Gain versus overhead

# Outline

## Motivation

JavaScript

Polyhedral Model And JavaScript

## JavaScriptCore

## Challenges And Solutions

Detection Of SCoPs

Parallel Speculation Failure

Gain vs Overhead

## Preliminary Results

# JavaScript Is Dynamic → We Need An Engine

```
f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
        for (var j = 0; j < height; i++) {  
            var v = img[i*width + j];  
            v = v + 41;  
            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

# JavaScript Is Dynamic → We Need An Engine

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- Types are dynamic

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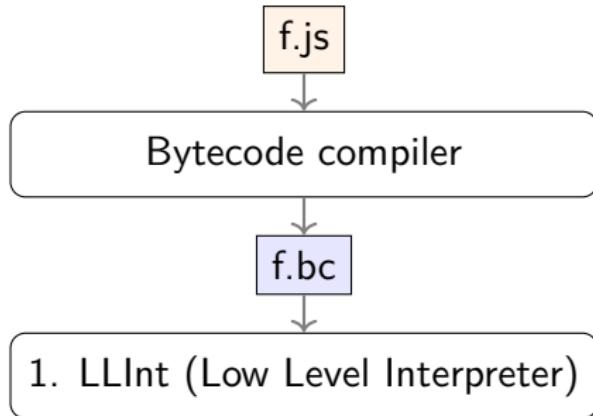
- Types are dynamic
- Arrays are dynamic

# JavaScript Is Dynamic → We Need An Engine

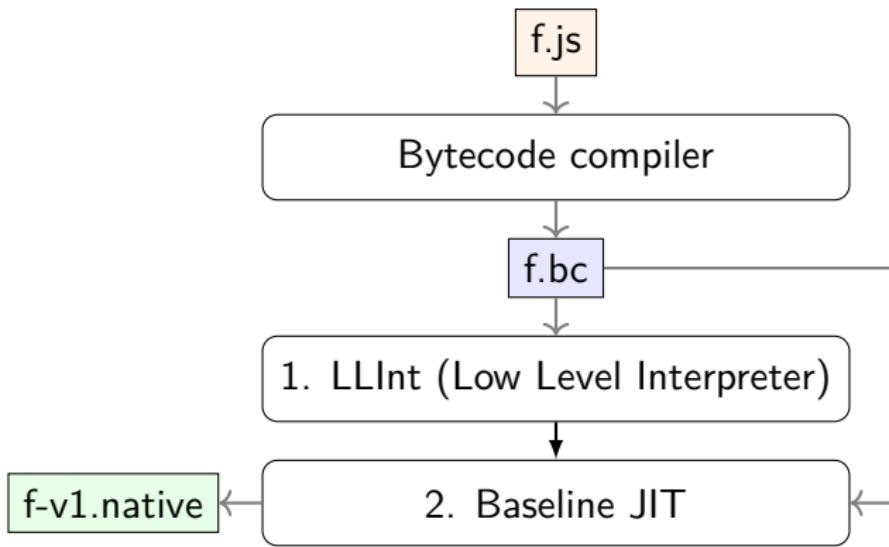
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            v = v + 41;  
            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

- Types are dynamic
- Arrays are dynamic
- **Only** double precision floating point numbers

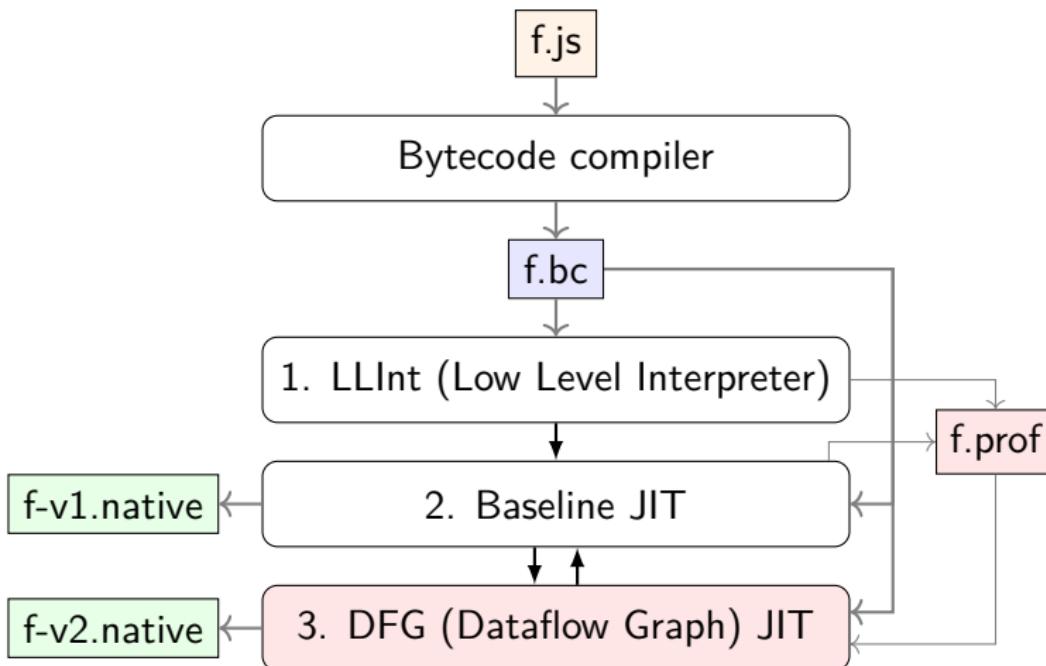
## 4 Layers



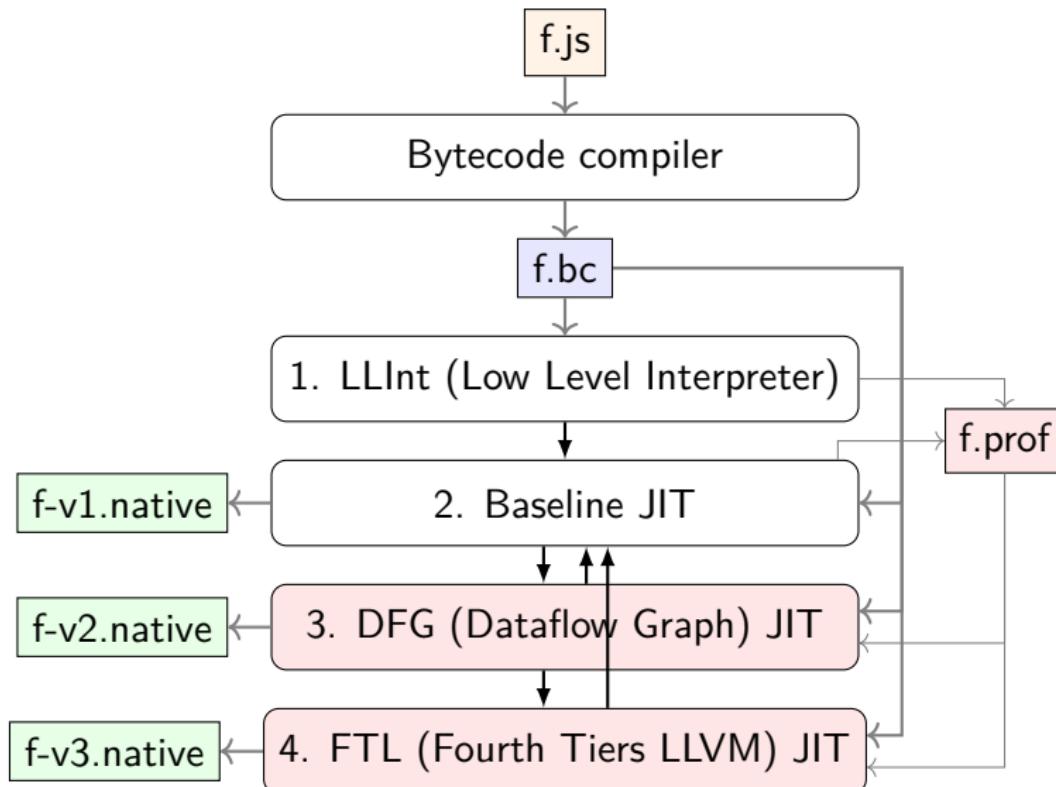
## 4 Layers



## 4 Layers



## 4 Layers



# 1. LLInt

```
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Source

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        }  
    }  
}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
    op_mul v 2 v;  
    ...  
    ...  
}
```

Bytecode

# 1. LLInt

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f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
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            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
    op_mul v 2 v;  
    ...  
    ...  
}
```

Bytecode

```
while(i = next_instruction()) {  
    switch(i->opcode) {  
        case add:  
            switch (type_pair(i->operand1->type(), i->operand2->type())):  
                case number_number:  
                    i->dest = add(i->operand1, i->operand2);  
                case object_number:  
                    i->dest = add(i->operand1->as_number(), i->operand2);  
                    ...  
        case mul: ...  
    }  
}
```

Interpreter

## 2. Baseline JIT

```
f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
        for (var j = 0; j < height; i++) {  
            var v = img[i*width + j];  
            v = v + 41;  
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        }  
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}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
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    ...  
    ...  
}
```

Bytecode

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f(img, width, height) {  
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            img[i*width + j] = v;  
        }  
    }  
}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
    op_mul v 2 v;  
    ...  
    ...  
}
```

Bytecode

```
...  
switch (type_pair(v, 41)): {  
    case number_number:  
        v = add(v, 41);  
    case object_number:  
        ...  
}
```

Native code

```
switch (type_pair(v, 2)): {  
    case number_number:  
        v = mul(v->as_number() * 2);  
    case object_number:  
        ...  
}
```

### 3. DFG JIT

```
f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
        for (var j = 0; j < height; i++) {  
            var v = img[i*width + j];  
            v = v + 41;  
            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
    op_mul v 2 v;  
    ...  
    ...  
}  
} Bytecode
```

### 3. DFG JIT

```
f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
        for (var j = 0; j < height; i++) {  
            var v = img[i*width + j];  
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            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
    op_mul v 2 v;  
    ...  
    ...  
}
```

Bytecode

```
if (!is_32int_array(img)) { return to Baseline JIT; }  
...  
res = v->as_32int() + 41;  
if (overflow(res)) { return to Baseline JIT; }  
v = res;  
res = v->as_32int() * 2;  
if (overflow(res)) { return to Baseline JIT; }  
v = res;  
...
```

DFG IR - Typed

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Source

```
f(img, width, height) {  
    ...  
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if (overflow(res)) { return to Baseline JIT; }  
v = res;  
res = v->as_32int() * 2;  
if (overflow(res)) { return to Baseline JIT; }  
v = res;  
...
```

DFG IR - Typed

Homemade  
optim and  
backend

```
...  
...
```

Native code

## 4. FTL JIT

```
f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
        for (var j = 0; j < height; i++) {  
            var v = img[i*width + j];  
            v = v + 41;  
            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

Source

```
f(img, width, height) {  
    ...  
    ...  
    op_add v 41 v;  
    op_mul v 2 v;  
    ...  
    ...  
}  
} Bytecode
```

## 4. FTL JIT

```
f(img, width, height) {  
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}
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Source

```
f(img, width, height) {  
    ...  
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}
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Bytecode

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if (!is_32int_array(img)) { return to Baseline JIT; }  
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res = v->as_32int() + 41;  
if (overflow(res)) { return to Baseline JIT; }  
v = res;  
res = v->as_32int() * 2;  
if (overflow(res)) { return to Baseline JIT; }  
v = res;  
...
```

LLVM IR - Typed

## 4. FTL JIT

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f(img, width, height) {  
    for (var i = 0; i < width; i++) {  
        for (var j = 0; j < height; i++) {  
            var v = img[i*width + j];  
            v = v + 41;  
            v = v * 2;  
            img[i*width + j] = v;  
        }  
    }  
}
```

Source

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f(img, width, height) {  
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if (overflow(res)) { return to Baseline JIT; }  
v = res;  
...
```

LLVM IR - Typed

LLVM  
optim and  
backend

```
...  
...
```

Native code

# Proposal

## Add polyhedral optimization in JavaScriptCore

- In the last layer - FTL
  - Dynamism has been removed
- Polly
  - Polyhedral optimizer for LLVM
  - Transform LLVM IR to optimized LLVM IR

# Outline

## Motivation

- JavaScript

- Polyhedral Model And JavaScript

## JavaScriptCore

## Challenges And Solutions

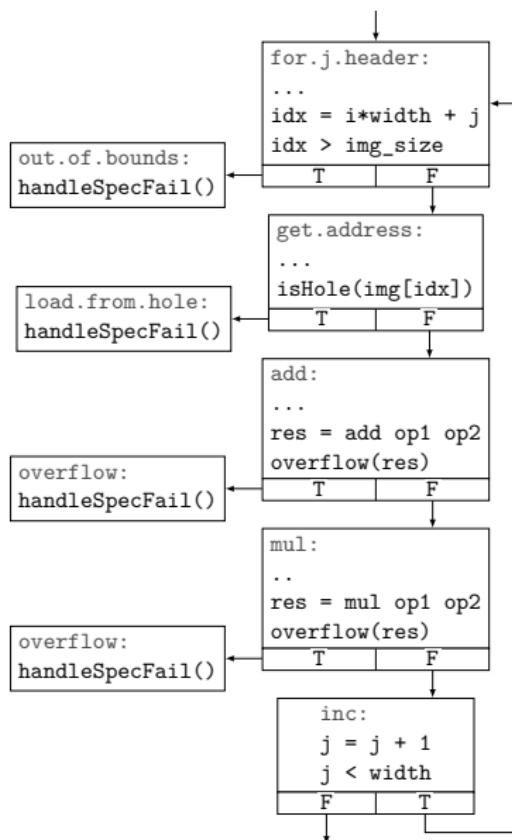
- Detection Of SCoPs

- Parallel Speculation Failure

- Gain vs Overhead

## Preliminary Results

# SESE Regions: Problem - Apply To All Engines

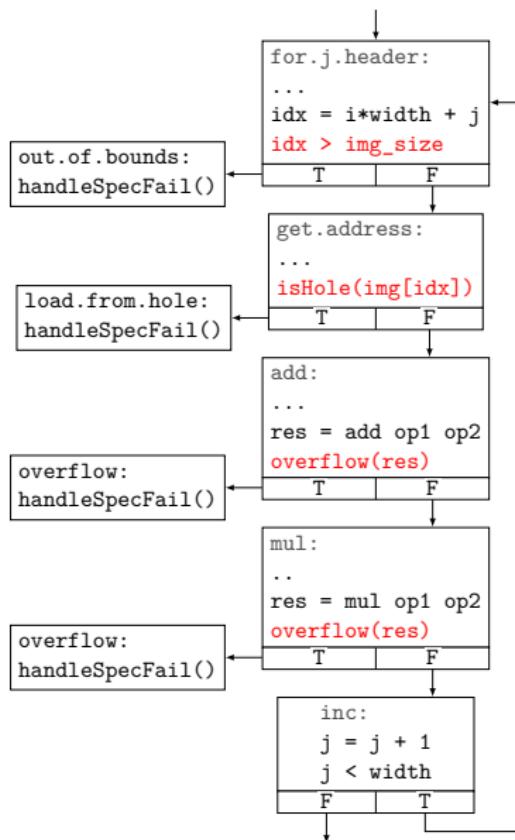


```

f(img, width, height) {
    for (var i = 0; i < width; i++) {
        for (var j = 0; j < height; i++) {
            var v = img[i*width + j]
            v = v + 41;
            v = v * 2;
            img[i*width + j] = v;
        }
    }
}
  
```

**Source**

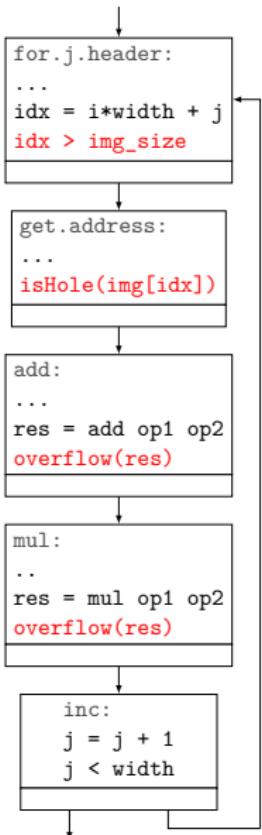
# SESE Regions: Solution



Step 1 (*in our implementation*)

- Tag instructions branching to exit blocks

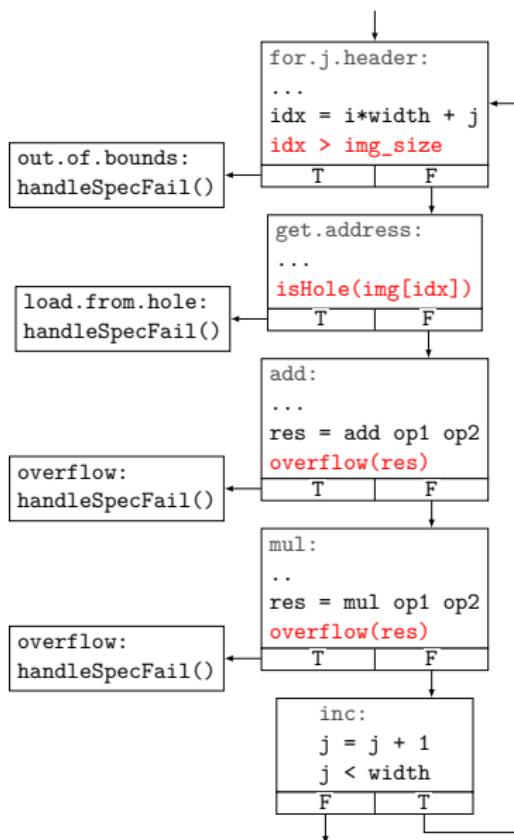
# SESE Regions: Solution



## Step 1 (*in our implementation*)

- Tag instructions branching to exit blocks
- Remove exit blocks
- Perform polyhedral optimization

# SESE Regions: Solution



## Step 1 (*in our implementation*)

- Tag instructions branching to exit blocks
- Remove exit blocks
- Perform polyhedral optimization

## Step 2 (*not yet completed*)

- Add back exit blocks from tagged instructions

# Two Dimensional Arrays And Arrays Of Objects

Problem - *Apply to all engines*

```
t[i].foo = 17;  
t[i][j] = 17;
```

→ 2 memory accesses

# Two Dimensional Arrays And Arrays Of Objects

## Problem - *Apply to all engines*

```
t[i].foo = 17;  
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```

→ 2 memory accesses

## Solutions

- Do not handle them (*in our implementation*)
- Inspector / executor
- Modify language and object allocator

# Handling `inttoptr` And `sext` Instructions

## Problem - *Specific to JavaScriptCore and Polly*

- Polly is designed for LLVM IR coming from frontends for static languages
- `intoptr` used by the runtime for known addresses
- `sext` used by the runtime for values representation (NaN boxing)
- `inttoptr` and `sext` instructions make SCoP detection fail

## Solution (*in our implementation*)

- Modify Polly to support these instructions

# Parallel Speculation

## Problem - *Apply to all engines*

- Speculation failure in one thread
- Other threads may have performed wrong computation

## Solutions

- Ignore speculation failure (*in our implementation*) - **Wrong!**
- Only optimize idempotent regions
- Save and rollback

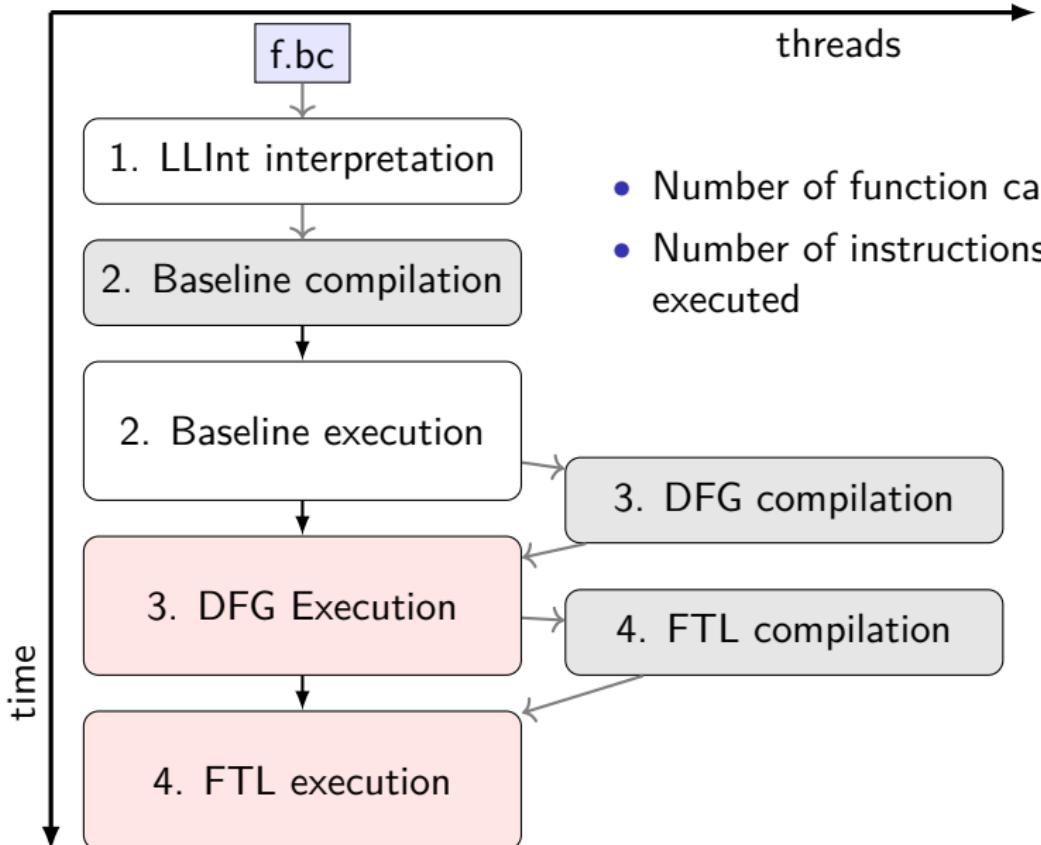
# Gain vs Overhead - Problem - *Apply To All Engines*

Execution time of optimized version + Polly optimization time

<

Execution time of original version

# Solution - JavaScriptCore Cost Model + Parallel Compil.



# Outline

## Motivation

- JavaScript

- Polyhedral Model And JavaScript

## JavaScriptCore

## Challenges And Solutions

- Detection Of SCoPs

- Parallel Speculation Failure

- Gain vs Overhead

## Preliminary Results

# Matrix Multiply - Transformations

```
matmul(left, right, res, left_nblines, left_nbcols, right_nbcols) {
    for (var i = 0; i < left_nblines; i++) {
        for (var j = 0; j < left_nbcols; j++) {
            var idx_left = i * left_nbcols + j;
            for (var k = 0; k < right_nbcols; k++) {
                var idx_res = i*right_nbcols + k;
                var idx_right = j*right_nbcols + k;
                res[idx_res] = res[idx_res] + left[idx_left]*right[idx_right];
            }
        }
    }
}
```

# Matrix Multiply - Transformations

```
matmul(left, right, res, left_nblines, left_nbcols, right_nbcols) {
    for (var i = 0; i < left_nblines; i++) {
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            var idx_left = i * left_nbcols + j;
            for (var k = 0; k < right_nbcols; k++) {
                var idx_res = i*right_nbcols + k;
                var idx_right = j*right_nbcols + k;
                res[idx_res] = res[idx_res] + left[idx_left]*right[idx_right];
            }
        }
    }
}
```

```
if(alias test ok){
    #pragma omp parallel for
    for(c0 = 0; c0 <= floord(p2-1, 32); c0 += 1)
        for(c1 = 0; c1 <= floord(p1-1, 32); c1 += 1)
            for(c2 = 0; c2 <= floord(p0-1, 32); c2 += 1) {
                for(c3 = 0; c3 <= min(31, p2-32*c0-1); c3 += 1)
                    for(c4 = 0; c4 <= min(31, p1-32*c1-1); c4 += 1)
                        for(c5 = 0; c5 <= min(31, p0-32*c2-1); c5 += 1)
                            Stmt_68(32*c0+c3, 32*c2+c5, 32*c1+c4);
    } else { original code version }
}
```

# Matrix Multiply - Performances

## Setup

- Intel Xeon W3520 with 4 cores running Linux 4.4.0
- LLVM and Polly 4.0.0 with --parallel
- `matmul` function called twice
- Right matrix size is 3000x300

## Results

Size of left matrix	Execution time without Polly (s)	Execution time with Polly (s)	Speedup
50x3000	0.08	0.06	1.33
500x3000	0.85	0.22	3.86
2000x3000	3.3	0.87	3.79

# Conclusion

Polyhedral optimization for JavaScript is almost there

- Speedups shown on a matrix multiply kernel

Ongoing work

- Complete implementation
- Evaluate speedups on real applications
- Study how to enrich JavaScriptCore profiling to help polyhedral optimizer

## An Other Conclusion

JavaScript is a dynamic language. To efficiently execute it, dynamic features must not be used too often. Engineers at Google writing the V8 engine even recommend<sup>1</sup> to:

*“Write code that looks like statically typed”*

---

<sup>1</sup>Franziska Hinkelmann at JS Conf EU 2017

# An Other Conclusion

JavaScript is a dynamic language. To efficiently execute it, dynamic features must not be used too often. Engineers at Google writing the V8 engine even recommend<sup>1</sup> to:

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**WebAssembly**

---

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# BACKUP

# Tiers-Up In JavaScriptCore

## LLInt → Baseline

- Same program state (variables values) representation
- OSR entry = jump (nothing to do)
- At any bytecode instruction

## Baseline → DFG

- Different program state representations
- OSR entry = jump + copy of the state
- Only at function entry and loop header

## DFG → FTL

- Different program state representations
- OSR entry = jump + copy of the state
- Only at function entry and loop header
  - Two code versions

# NaN-Boxing In JavaScriptCore

	Tag	Payload	
IEE-754 doubles	{	0 0 0 0   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F 8   F F F F   F F F F   F F F F	}
			JavaScriptCore doubles
IEE-754 NaN	{	F F F 9   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F D   F F F F   F F F F   F F F F	}
			JavaScriptCore NaN
	{	F F F E   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F E   0 0 0 0   F F F F   F F F F	}
			32-bits integers
		.....	Not used
	{	F F F F   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F F   F F F F   F F F F   F F F F	}
			JavaScriptCore pointers

- **Direct** manipulation of doubles
- **Indirect** manipulation (masking) of pointers
- **Indirecte** manipulation (masking) 32 bits integers

# NaN-Boxing In JavaScriptCore

	Tag	Payload	
IEEE-754 doubles	{	0 0 0 0   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F 8   F F F F   F F F F   F F F F	}
	{	F F F 9   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F F   F F F F   F F F F   F F F F	}
II N	Favor pointers	<ul style="list-style-type: none"> <li><math>+2^{48} = 0001 0000 0000 0000 = 281474976710656</math></li> </ul>	ptCore
	{	FFFF   0 0 0 0   0 0 0 0   0 0 0 0 ..... FFFF   F F F F   F F F F   F F F F	}
			pointers

- **Direct** manipulation of doubles
- **Indirect** manipulation (masking) of pointers
- **Indirecte** manipulation (masking) 32 bits integers

# NaN-Boxing In JavaScriptCore

	Tag	Payload	
IEE-754 doubles	{	0 0 0 1   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F 9   F F F F   F F F F   F F F F	JavaScriptCore doubles
IEE-754 NaN	{	F F F A   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F E   F F F F   F F F F   F F F F	JavaScriptCore NaN
	{	F F F F   0 0 0 0   0 0 0 0   0 0 0 0 ..... F F F F   0 0 0 0   F F F F   F F F F	32-bits integers
		.....	Not used
	{	0 0 0 0   0 0 0 0   0 0 0 0   0 0 0 0 ..... 0 0 0 0   F F F F   F F F F   F F F F	JavaScriptCore pointers

- **Direct** manipulation of pointers
- **Indirect** manipulation (substraction) of doubles
- **Indirect** manipulation (masking) of 32 bits integers

# NaN-Boxing In JavaScriptCore

- DoubleEncodeOffset
  - Favor pointers
  - $0001|0000|0000|0000 = +281474976710656 = +2^{48}$
- TagTypeNumber
  - If all bits in the mask are set, this indicates an integer
  - If any but not all are set this value is a double.
  - $FFFF|0000|0000|0000 = -281474976710656$
- TagMask = TagTypeNumber | TagBitTypeOther
  - Used to check for all types of immediate values
  - Either number or other immediate (bool, null, undefined)
  - $FFFF|0000|0000|0002 = -281474976710654$
- add DoubleEncodeOffset  $\equiv$  sub TagTypeNumber

# JavaScriptCore History

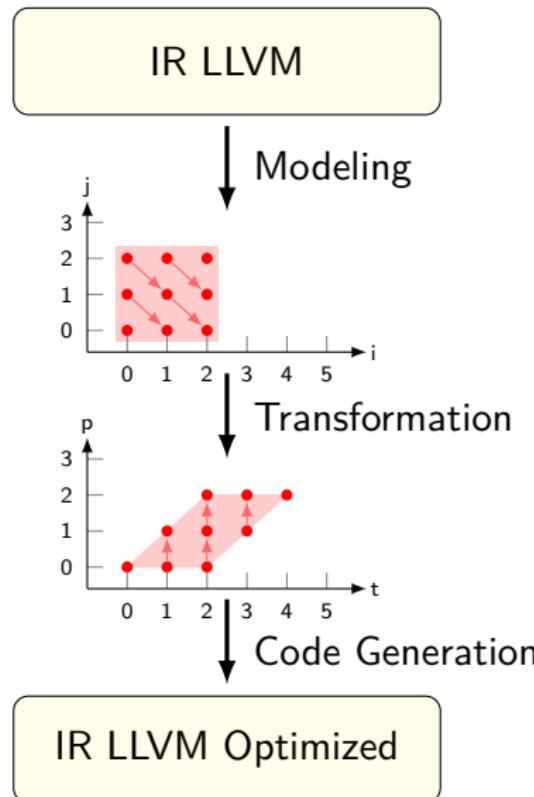
## Developed by Apple

- Included in WebKit, LGPL license
- Before 2008 - SquirrelFish interpreter
- 2008 - 2014 - Nitro JIT
- 2014 - FTL JIT based on LLVM
- Mid 2016 - FTL replaced by B3, home made JIT and backend

## Big project

LOC per language:	
cpp:	289 435 (89.38%)
ansic:	11 254 (3.48%)
ruby:	9 925 (3.06%)
python:	6 195 (1.91%)
asm:	4 982 (1.54%)
perl:	2 013 (0.62%)
sh:	24 (0.01%)
LOC Total:	323 828

# Polly: Polyhedral Optimization for LLVM



# Detection Of Affine Accesses To Arrays

Problem - *Specific to JavaScriptCore and Polly*

```
t[index] = 17;
```

```
%offset = shl i64 %index, 3
%cell_as_int = add i64 %base_as_int, %offset
%cell_ptr = inttoptr i64 %cell_as_int to i64*
store i64 %boxed_17, i64* %cell_ptr
```

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Problem - *Specific to JavaScriptCore and Polly*

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```

Solution (*in our implementation*)

```
%base_ptr = inttoptr i64 %base_as_int
            to [1000 x i64]*
%cell_ptr = getelementptr [1000 x i64],
            [1000 x i64]* %base_ptr,
            i32 0, i32 %index
store i64 %boxed_17, i64* %cell_ptr
```